

HIGH-SPEED 2.5V 512/256K x 18 ASYNCHRONOUS DUAL-PORT STATIC RAM WITH 3.3V OR 2.5V INTERFACE

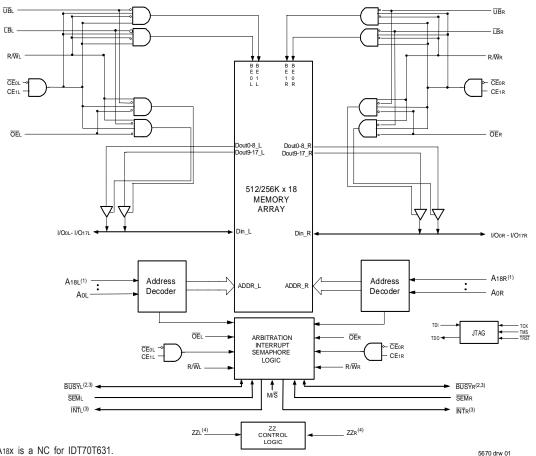
PRELIMINARY IDT70T633/1S

Features

- True Dual-Port memory cells which allow simultaneous access of the same memory location
- High-speed access
 - Commercial: 10/12/15ns (max.)
 - Industrial: 12ns (max.)
- Dual chip enables allow for depth expansion without external logic
- IDT70T633/1 easily expands data bus width to 36 bits or more using the Master/Slave select when cascading more than one device
- ◆ M/S = VIH for BUSY output flag on Master, $M/\overline{S} = VIL$ for \overline{BUSY} input on Slave
- Busy and Interrupt Flags
- Full hardware support of semaphore signaling between ports on-chip

- On-chip port arbitration logic
- Fully asynchronous operation from either port
- Separate byte controls for multiplexed bus and bus matching compatibility
- Sleep Mode Inputs on both ports
- Supports JTAG features compliant to IEEE 1149.1 in BGA-208 and BGA-256 packages
- Single 2.5V (±100mV) power supply for core
- LVTTL-compatible, selectable 3.3V (±150mV)/2.5V (±100mV) power supply for I/Os and control signals on each port
- Available in a 256-ball Ball Grid Array, 144-pin Thin Quad Flatpack and 208-ball fine pitch Ball Grid Array
- Industrial temperature range (-40°C to +85°C) is available for selected speeds

Functional Block Diagram



NOTES:

- Address A₁₈x is a NC for IDT70T631.
- BUSY is an input as a Slave (M/S=VIL) and an output when it is a Master (M/S=VIH).
- BUSY and INT are non-tri-state totem-pole outputs (push-pull).
- 4. The sleep mode pin shuts off all dynamic inputs, except JTAG inputs, when asserted. OPTx, BUSYx, INTx, M/S and the sleep mode pins themselves (ZZx) are not affected during sleep mode.

APRIL 2003

Description

The IDT70T633/1 is a high-speed 512/256K x 18 Asynchronous Dual-Port Static RAM. The IDT70T633/1 is designed to be used as a stand-alone 9216/4608K-bit Dual-Port RAM or as a combination MAS-TER/SLAVE Dual-Port RAM for 36-bit-or-more word system. Using the IDT MASTER/SLAVE Dual-Port RAM approach in 36-bit or wider memory system applications results in full-speed, error-free operation without the need for additional discrete logic.

This device provides two independent ports with separate control,

address, and I/O pins that permit independent, asynchronous access for reads or writes to any location in memory. An automatic power down feature controlled by the chip enables (either $\overline{CE}0$ or CE1) permit the on-chip circuitry of each port to enter a very low standby power mode.

The 70T633/1 can support an operating voltage of either 3.3V or 2.5V on one or both ports, controlled by the OPT pins. The power supply for the core of the device (VDD) remains at 2.5V.

Pin Configuration^(1,2,3)

70T633/1BC BC-256(5,6)

256-Pin BGA Top View

03/13/03

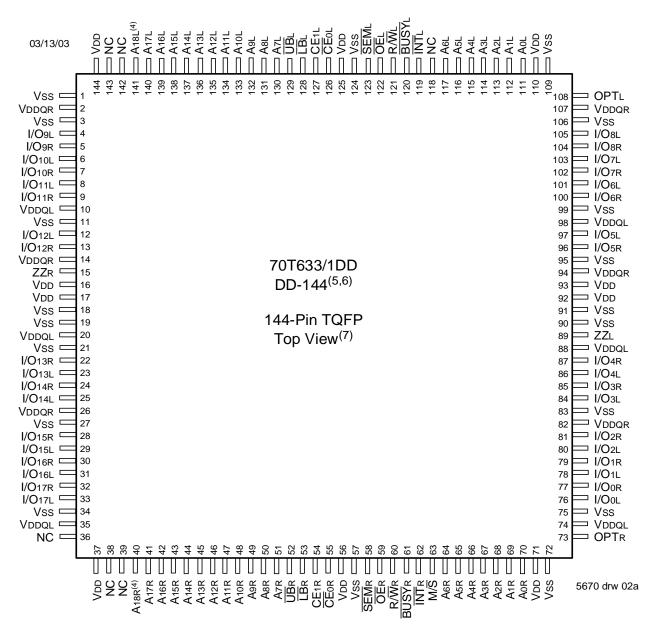
A1	A2	A3	A4	A5	A6	A7	A8	A9	A10	A11	A12	A13	A14	A15	A16
NC	TDI	NC	A17L	A14L	A11L	A8L	NC	CE1L	OE L	INTL	A 5L	A2L	A0L	NC	NC
B1	B2	B3	B4	B5	B6	B7	B8	B9	B10	B11	B12	B13	B14	B15	B16
NC	NC	TDO	A18L ⁽⁴⁾	A15L	A12L	A9L	UBL	CE0L	R/WL	NC	A4L	A1L	NC	NC	NC
C1	C2	C3	C4	C5	C6	C7	C8	C9	C10	C11	C12	C13	C14	C15	C16
NC	I/O9L	Vss	A16L	A13L	A10L	A7L	NC	LBL	SEML	BUSYL	A6L	A3L	OPTL	NC	I/O8L
D1	D2	D3	D4	d5	D6	d7	d8	d9	d10	d11	D12	D13	D14	D15	D16
NC	I/O9R	NC	Vdd	Vddql	Vddql	Vddqr	Vddqr	Vddql	Vddql	Vddqr	Vddqr	Vdd	NC	NC	I/O8R
E1	E2	E3	E4	E5	E6	E7	E8	E9	E10	E11	E12	E13	E14	E15	E16
I/O10R	I/O10L	NC	Vddql	Vdd	Vdd	Vss	Vss	Vss	Vss	VDD	VDD	Vddqr	NC	I/O7L	I/ O 7R
F1	F2	F3	F4	F5	F6	F7	F8	F9	F10	F11	F12	F13	F14	F15	F16
I/O11L	NC	I/O11R	Vddql	Vdd	NC	Vss	Vss	Vss	Vss	Vss	Vdd	Vddqr	I/O6R	NC	I/O6L
G1	G2	G3	g4	G5	G6	G7	G8	_{G9}	G10	G11	G12	G13	G14	G15	G16
NC	NC	I/O12L	Vddqr	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddql	I/O5L	NC	NC
H1	H2	H3	h4	H5	H6	H7	H8	H9	H10	H11	H12	H13	H14	H15	H16
NC	I/O12R	NC	Vddqr	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddql	NC	NC	I/O5R
J1	J2	J3	J4	J5	J6	J7	J8	^{J9}	J10	J11	J12	J13	J14	J15	J16
I/O13L	I/O14R	I/O13R	Vddql	ZZR	Vss	Vss	Vss	Vss	Vss	Vss	ZZ L	Vddqr	I/O4R	I/O3R	I/O4L
K1	K2	K3	K4	K5	K6	K7	K8	K9	K10	K11	K12	K13	K14	K15	K16
NC	NC	I/O14L	Vddql	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddqr	NC	NC	I/O3L
L1	L2	L3	L4	L5	L6	L7	L8	L9	L10	L11	L12	L13	L14	L15	L16
I/O15L	NC	I/O15R	Vddqr	Vdd	NC	Vss	Vss	Vss	Vss	Vss	Vdd	Vddql	I/O2L	NC	I/O2R
M1	M2	M3	m4	M5	M6	M7	M8	м9	M10	M11	M12	M13	M14	M15	M16
I/O16R	I/O16L	NC	Vddqr	Vdd	Vdd	Vss	Vss	Vss	Vss	Vdd	Vdd	Vddql	I/O1R	I/O1L	NC
N1	N2	N3	N4	N5	N6	n7	n8	n9	N10	N11	N12	N13	N14	N15	N16
NC	I/O17R	NC	Vdd	Vddqr	Vddqr	Vddql	Vddql	Vddqr	Vddqr	Vddql	Vddql	Vdd	NC	I/O0R	NC
P1	P2	P3	P4	P5	P6	P7	P8	P9	P10	P11	P12	P13	P14	P15	P16
NC	I/O17L	TMS	A16R	A13R	A10R	A 7R	NC	LBr	SEMR	BUSYR	A 6R	A 3R	NC	NC	I/OoL
R1	R2	R3	R4	R5	R6	R7	R8	R9	R10	R11	R12	R13	R14	R15	R16
NC	NC	TRST	A 18R ⁽⁴⁾	A15R	A12R	A 9R	UBr	CEor	R/WR	M/S	A 4R	A1R	OPTR	NC	NC
T1	T2	T3	T4	T5	T6	T7	T8	T9	T10	T11	T12	T13	T14	T15	T16
NC	TCK	NC	A17R	A 14R	A 11R	A8R	NC	CE1R	OE R	INT _R	A 5R	A 2R	Aor	NC	NC

NOTES:

5670 drw 02c

- 1. All $\ensuremath{\mathsf{VDD}}$ pins must be connected to 2.5V power supply.
- 2. All VDDQ pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (2.5V), and 2.5V if OPT pin for that port is set to Vss (0V).
- 3. All Vss pins must be connected to ground supply.
- 4. A18x is a NC for IDT70T631.
- 5. Package body is approximately 17mm x 17mm x 1.4mm, with 1.0mm ball-pitch.
- 6. This package code is used to reference the package diagram.

Pin Configurations^(1,2,3,8) (con't.)



- 1. All VDD pins must be connected to 2.5V power supply.
- 2. All VDDQ pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (2.5V), and 2.5V if OPT pin for that port is set to Vss (0V).
- 3. All Vss pins must be connected to ground.
- 4. A18x is a NC for IDT70T631.
- 5. Package body is approximately 20mm x 20mm x 1.4mm.
- 6. This package code is used to reference the package diagram.
- 7. This text does not indicate orientation of the actual part-marking.
- 8. Due to the restricted number of pins, JTAG is not supported in the DD-144 package.

Pin Configurations^(1,2,3)(con't.)

03/12	2/03																	
	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	
Α	I/O ₉ L	NC	Vss	TDO	NC	A _{16L}	A ₁₂ L	AsL	NC	VDD	SEML	ĪNT∟	A4L	Aol	OPT∟	NC	Vss	Α
В	NC	Vss	NC	TDI	A17L	A13L	A9L	NC	<u>CE</u> ₀L	Vss	BUSŢL	A ₅ L	A ₁ L	Vss	VDDQR	I/O ₈ L	NC	В
С	VDDQL	I/O ₉ R	VDDQR	V _{DD}	A _{18L} ⁽⁴⁾	A14L	A _{10L}	ŪB∟	CE ₁ L	Vss	R/WL	A ₆ L	A ₂ L	VDD	I/O _{8R}	NC	Vss	С
D	NC	Vss	I/O _{10L}	NC	A ₁₅ L	A ₁₁ L	A7L	ΪΒι	VDD	ŌĒL	NC	Азь	VDD	NC	VDDQL	I/O7L	I/O7R	D
Е	I/O11L	NC	VDDQR	I/O _{10R}										I/O ₆ L	NC	Vss	NC	Ε
F	VDDQL	I/O11R	NC	Vss										Vss	I/O ₆ R	NC	VDDQR	F
G	NC	Vss	I/O _{12L}	NC										NC	VDDQL	I/O ₅ L	NC	G
Н	VDD	NC	VDDQR	I/O _{12R}					T633 F-208					VDD	NC	Vss	I/O ₅ R	Н
J	VDDQL	VDD	Vss	ZZR		208-Ball BGA								ZZι	VDD	Vss	VDDQR	J
K	I/O _{14R}	Vss	I/O13R	Vss					p Vie		•			I/O3R	VDDQL	I/O ₄ R	Vss	K
L	NC	I/O14L	VDDQR	I/O13L										NC	I/O3L	Vss	I/O ₄ L	L
М	VDDQL	NC	I/O _{15R}	Vss										Vss	NC	I/O ₂ R	VDDQR	M
Ν	NC	Vss	NC	I/O ₁₅ L										I/O1R	VDDQL	NC	I/O ₂ L	N
Р	I/O _{16R}	I/O _{16L}	VDDQR	NC	TRST	A _{16R}	A ₁₂ R	Asr	NC	Voo	SEMR	ĪÑ₹r	A ₄ R	NC	I/O1L	Vss	NC	Р
R	Vss	NC	I/O17R	тск	A17R	A _{13R}	A ₉ R	NC	<u>CE</u> or	Vss	BUSYR	A ₅ R	A ₁ R	Vss	VDDQL	I/Oor	VDDQR	R
Т	NC	I/O17L	VDDQL	TMS	A _{18R} ⁽⁴⁾	A _{14R}	A10R	ŪBR	CE _{1R}	Vss	R/W̄R	A ₆ R	A ₂ R	Vss	NC	Vss	NC	Т
U	Vss	NC	V _{DD}	NC	A _{15R}	A _{11R}	A7R	ŪBR	VDD	ŌĒR	M/S	Азп	Aor	VDD	OPTR	NC	I/OoL	U

5670 drw 02b

- 1. All VDD pins must be connected to 2.5V power supply.
- 2. All VDDQ pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (2.5V), and 2.5V if OPT pin for that port is set to Vss (0V).
- 3. All Vss pins must be connected to ground.
- 4. A₁₈x is a NC for IDT70T631.
- 5. Package body is approximately 15mm x 15mm x 1.4mm with 0.8mm ball pitch.
- 6. This package code is used to reference the package diagram.
- 7. This text does not indicate orientation of the actual part-marking.

Pin Names

Right Port	Names			
CEOR, CE1R	Chip Enables (Input)			
R/WR	Read/Write Enable (Input)			
ŌĒr	Output Enable (Input)			
A0R - A18R ⁽¹⁾	Address (Input)			
I/O0R - I/O17R	Data Input/Output			
<u>SEM</u> _R	Semaphore Enable (Input)			
ĪNTR	Interrupt Flag (Output)			
BUSYR	Busy Flag (Output)			
UB R	Upper Byte Select (Input)			
<u>IB</u> _R	Lower Byte Select (Input)			
VDDQR	Power (I/O Bus) (3.3V or 2.5V) ⁽²⁾ (Input)			
OPTr	Option for selecting VDDqx ^(2,3) (Input)			
ZZ R	Sleep Mode Pin ⁽⁴⁾ (Input)			
M/S	Master or Slave Select (Input)			
VDD	Power (2.5V) ⁽²⁾ (Input)			
Vss	Ground (0V) (Input)			
TDI	Test Data Input			
TDO	Test Data Output			
ГСК	Test Logic Clock (10MHz) (Input)			
ΓMS	Test Mode Select (Input)			
RST	Reset (Initialize TAP Controller) (Input)			
	CEOR, CE1R R/WR OER AOR - A18R(1) I/OOR - I/O17R SEMR INTR BUSYR UBR LBR VDDQR OPTR			

5670 tbl 01

- 1. Address A₁₈x is a NC for IDT70T631.
- VDD, OPTx, and VDDQx must be set to appropriate operating levels prior to applying inputs on I/Ox.
- 3. OPTx selects the operating voltage levels for the I/Os and controls on that port. If OPTx is set to Vpp (2.5V), then that port's I/Os and controls will operate at 3.3V levels and Vppox must be supplied at 3.3V. If OPTx is set to Vss (0V), then that port's I/Os and controls will operate at 2.5V levels and Vppox must be supplied at 2.5V. The OPT pins are independent of one another—both ports can operate at 3.3V levels, both can operate at 2.5V levels, or either can operate at 3.3V with the other at 2.5V.
- The sleep mode pin shuts off all dynamic inputs, except JTAG inputs, when asserted. OPTx, BUSYx, INTx, M/S and the sleep mode pins themselves (ZZx) are not affected during sleep mode. JTAG cannot be tested in sleep mode.
- BUSY is an input as a Slave (M/S=VIL) and an output when it is a Master (M/S=VIH).

Truth Table I—Read/Write and Enable Control⁽¹⁾

ŌĒ	SEM	<u>CE</u> ₀	CE ₁	ŪB	ΙB	R/W	7 Z	Upper Byte	Lower Byte	MODE
UE	SEIVI	CEO	CE1	UB	LB	FC/ VV	11	1/09-1/	1/00-8	MODE
Х	Н	Н	X	Χ	Х	Х	L	High-Z	High-Z	Deselected–Power Down
Х	Н	Χ	ш	Х	Х	Х	ш	High-Z	High-Z	Deselected-Power Down
Х	Н	Ш	Ι	Ι	Η	Х	Ш	High-Z	High-Z	Both Bytes Deselected
Х	Н	L	Н	Н	L	L	L	High-Z	Din	Write to Lower Byte
Х	Н	L	Η	L	Н	L	L	Din	High-Z	Write to Upper Byte
Х	Н	L	Н	L	L	L	L	Din	Din	Write to Both Bytes
L	Н	L	Η	Н	L	Н	L	High-Z	Dout	Read Lower Byte
L	Н	L	Н	L	Н	Н	L	D оит	High-Z	Read Upper Byte
L	Н	L	Н	L	L	Н	L	D оит	Dout	Read Both Bytes
Н	Н	L	Н	L	L	Х	L	High-Z	High-Z	Outputs Disabled
Х	Х	Χ	Х	Х	Х	Х	Н	High-Z	High-Z	High-Z Sleep Mode

5670 tbl 02

NOTE:

1. "H" = VIH, "L" = VIL, "X" = Don't Care.

Truth Table II - Semaphore Read/Write Control(1)

	Inputs ⁽¹⁾					Outputs		
CE(2)	R/W	ŌĒ	ŪB	ĪΒ	SEM	I/O1-17 I/O0		Mode
Н	Н	L	L	L	L	DATAout	DATAout	Read Data in Semaphore Flag ⁽³⁾
Н	↑	Χ	Χ	L	L	Х	DATAIN	Write I/Oo into Semaphore Flag
L	Χ	Χ	Χ	Χ	L			Not Allowed

- 1. There are eight semaphore flags written to I/Oo and read from all the I/Os (I/Oo-I/O17). These eight semaphore flags are addressed by Ao-A2.
- 2. $\overline{CE} = L$ occurs when $\overline{CE}_0 = V_{IL}$ and $CE_1 = V_{IH}$. $\overline{CE} = H$ when $\overline{CE}_0 = V_{IH}$ and/or $CE_1 = V_{IL}$.
- 3. Each byte is controlled by the respective \overline{UB} and \overline{LB} . To read data \overline{UB} and/or \overline{LB} = VIL.

NOTE:

Recommended Operating Temperature and Supply Voltage⁽¹⁾

Grade	Ambient Temperature	GND	VDD		
Commercial	0°C to +70°C	0V	2.5V <u>+</u> 100mV		
Industrial	-40°C to +85°C	0V	2.5V <u>+</u> 100mV		

5670 tbl 04

1. This is the parameter TA. This is the "instant on" case temperature.

Absolute Maximum Ratings(1)

Symbol	Rating	Commercial & Industrial	Unit
VTERM (VDD)	VDD Terminal Voltage with Respect to GND	-0.5 to 3.6	V
Vterm ⁽²⁾ (Vddq)	VDDQ Terminal Voltage with Respect to GND	-0.3 to VDDQ + 0.3	V
V _{TERM} ⁽²⁾ (INPUTS and I/O's)	Input and I/O Terminal Voltage with Respect to GND	-0.3 to VDDQ + 0.3	V
TBIAS ⁽³⁾	Temperature Under Bias	-55 to +125	°C
Тѕтс	Storage Temperature	-65 to +150	°C
TJN	Junction Temperature	+150	°C
IOUT(For VDDQ = 3.3V)	DC Output Current	50	mA
IOUT(For VDDQ = 2.5V)	DC Output Current	40	mA

5670 tbl 07

NOTES:

- Stresses greater than those listed under ABSOLUTE MAXIMUM RATINGS may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.
- This is a steady-state DC parameter that applies after the power supply has reached its nominal operating value. Power sequencing is not necessary; however, the voltage on any Input or I/O pin cannot exceed VDDQ during power supply ramp up.
- 3. Ambient Temperature under DC Bias. No AC Conditions. Chip Deselected.

Capacitance⁽¹⁾

TA = +25°C, F = 1.0MHz) TQFP ONLY

Symbol	Parameter	Conditions ⁽²⁾	Max.	Unit
Cin	Input Capacitance	VIN = 3dV	8	pF
Соит ⁽³⁾	Output Capacitance	Vout = 3dV	10.5	pF

5670 tbl 08

NOTES

- These parameters are determined by device characterization, but are not production tested.
- 3dV references the interpolated capacitance when the input and output switch from 0V to 3V or from 3V to 0V.
- 3. Cout also references CI/O.

Recommended DC Operating Conditions with VDDQ at 2.5V

Symbol	Parameter	Min.	Тур.	Max.	Unit
VDD	Core Supply Voltage	2.4	2.5	2.6	٧
VDDQ	I/O Supply Voltage ⁽³⁾	2.4	2.5	2.6	٧
Vss	Ground	0	0	0	٧
VIH	Input High Volltage (Address, Control & Data I/O Inputs) ⁽³⁾	1.7		VDDQ + 100mV ⁽²⁾	>
VIH	Input High Voltage - JTAG	1.7	_	VDD + 100mV ⁽²⁾	>
VIH	Input High Voltage - ZZ, OPT, M/S	VDD - 0.2V		VDD + 100mV ⁽²⁾	٧
VIL	Input Low Voltage	-0.3 ⁽¹⁾		0.7	٧
VIL	Input Low Voltage - ZZ, OPT, M/S	-0.3 ⁽¹⁾		0.2	٧

5670 tbl 05

NOTES:

- 1. VIL (min.) = -1.0V for pulse width less than tRC/2 or 5ns, whichever is less.
- 2. VIH (max.) = VDDQ + 1.0V for pulse width less than tRc/2 or 5ns, whichever is less
- To select operation at 2.5V levels on the I/Os and controls of a given port, the OPT pin for that port must be set to Vss (0V), and VDDQX for that port must be supplied as indicated above.

Recommended DC Operating Conditions with VDDQ at 3.3V

Symbol	Parameter	Min.	Тур.	Max.	Unit
VDD	Core Supply Voltage	2.4	2.5	2.6	٧
VDDQ	I/O Supply Voltage ⁽³⁾	3.15	3.3	3.45	٧
Vss	Ground	0	0	0	٧
VIH	Input High Voltage (Address, Control &Data I/O Inputs) ⁽³⁾	2.0	-	VDDQ + 150mV ⁽²⁾	٧
VIH	Input High Voltage - JTAG	1.7		VDD + 100mV ⁽²⁾	٧
VIH	Input High Voltage - ZZ, OPT, M/S	VDD - 0.2V	1	VDD + 100mV ⁽²⁾	٧
VIL	Input Low Voltage	-0.3 ⁽¹⁾		0.8	V
VL	Input Low Voltage - ZZ, OPT, M/S	-0.3 ⁽¹⁾		0.2	٧

5670 tbl 06

- 1. VIL (min.) = -1.0V for pulse width less than tRC/2 or 5ns, whichever is less.
- VIH (max.) = VDDQ + 1.0V for pulse width less than tRc/2 or 5ns, whichever is less.
- To select operation at 3.3V levels on the I/Os and controls of a given port, the OPT pin for that port must be set to VDD (2.5V), and VDDQX for that port must be supplied as indicated above.

5670 tbl 09

ight-speed 2.37 312/230K X to Asynchronous Buai-Fort Static NAM industrial and Commercial Temperature Nam

DC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range (VDD = 2.5V ± 100mV)

			70T63	33/1S	
Symbol	Parameter	Test Conditions	Min.	Max.	Unit
LI	Input Leakage Current ⁽¹⁾	VDDQ = Max., VIN = 0V to VDDQ	_	10	μΑ
LI	JTAG & ZZ Input Leakage Current ^(1,2)	V _{DD} = Max., V _{IN} = 0V to V _{DD}	_	<u>+</u> 30	μΑ
llo	Output Leakage Current(1,3)	CE0 = VIH or CE1 = VIL, VOUT = 0V to VDDQ	_	10	μΑ
Vol (3.3V)	Output Low Voltage ⁽¹⁾	IOL = +4mA, VDDQ = Min.	_	0.4	V
Vон (3.3V)	Output High Voltage(1)	IOH = -4mA, VDDQ = Min.	2.4	_	V
Vol (2.5V)	Output Low Voltage ⁽¹⁾	IOL = +2mA, VDDQ = Min.	_	0.4	V
Vон (2.5V)	Output High Voltage (1)	IOH = -2mA, VDDQ = Min.	2.0	_	V

NOTES:

- 1. VDDQ is selectable (3.3V/2.5V) via OPT pins. Refer to page 6 for details.
- 2. Applicable only for TMS, TDI and TRST inputs.
- 3. Outputs tested in tri-state mode.

DC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range (VDD = 2.5V ± 100mV)

70T633/1S10 70T633/1S12 70T633/1S15 Com'l Only Com'l Only Com'l & Ind Typ.(3) Typ.(3) Typ.(3) Symbol Parameter 1 4 1 **Test Condition** Version Max. Max. Max. Unit IDD Dynamic Operating CEL and CER= VIL, COM'L 500 650 415 540 330 430 mA Current (Both Outputs Disabled Ports Active) IND S 455 580 $f = fMAX^{(1)}$ ISB1⁽⁵⁾ Standby Current $\overline{CF}I = \overline{CF}R = VIH$ COM'L S 200 250 165 210 135 165 mΑ (Both Ports - TTL $f = fMAX^{(1)}$ Level Inputs) IND S 190 235 ISB2⁽⁵⁾ \overline{CE} "A" = VIL and \overline{CE} "B" = VIH⁽⁴⁾ Standby Current COM'L S 300 350 250 290 200 230 mΑ (One Port - TTL Active Port Outputs Disabled, Level Inputs) $f = fMAX^{(1)}$ IND S 275 315 ISB3 Full Standby Current Both Ports CEL and COM'L S 15 30 15 30 15 30 mΑ $\overline{\text{CE}}$ R \geq VDD - 0.2V, VIN \geq VDD - 0.2V (Both Ports - CMOS or VIN $\leq 0.2V$, f = $0^{(2)}$ IND Level Inputs) S 15 40 Full Standby Current ISB4(5) \overline{CE} "A" $\leq 0.2V$ and \overline{CE} "B" $\geq VDD - 0.2V^{(4)}$ COM'L S 300 350 250 290 200 230 mΑ (One Port - CMOS $VIN \ge VDD - 0.2V$ or $VIN \le 0.2V$, Active IND Level Inputs) Port, Outputs Disabled, f = fMAX(1) S 275 315 ΙZΖ Sleep Mode Current ZZL = ZZR = VIHmA COM'L S 15 30 15 30 15 30 (Both Ports - TTL $f = fMAX^{(1)}$ IND S 15 40 Level Inputs)

- 1. At f = fMAX, address and control lines (except Output Enable) are cycling at the maximum frequency read cycle of 1/tRc, using "AC TEST CONDITIONS".
- 2. f = 0 means no address or control lines change. Applies only to input at CMOS level standby.
- 3. VDD = 2.5V, $TA = 25^{\circ}C$ for Typ. values, and are not production tested. IDD DC(f=0) = 15mA (Typ).
- 4. $\overline{CE}x = V_{IL}$ means $\overline{CE}_{0x} = V_{IL}$ and $CE_{1x} = V_{IH}$
 - $\overline{CE}X = VIH \text{ means } \overline{CE}0X = VIH \text{ or } CE1X = VIL$
 - $\overline{\text{CE}}\text{x} \leq 0.2 \text{V}$ means $\overline{\text{CE}}\text{ox} \leq 0.2 \text{V}$ and $\text{CE}\text{1x} \geq \text{V}\text{DDQx}$ 0.2 V
 - $\overline{CE}x \ge VDDQx 0.2V$ means $\overline{CE}0x \ge VDDQx 0.2V$ or CE1x 0.2V
 - "X" represents "L" for left port or "R" for right port.
- 5. ISB1, ISB2 and ISB4 will all reach full standby levels (ISB3) on the appropriate port(s) if ZZL and /or ZZR = VIH.

AC Test Conditions (VDDQ = 3.3V/2.5V)

Input Pulse Levels	GND to 3.0V / GND to 2.5V
Input Rise/Fall Times	2ns Max.
Input Timing Reference Levels	1.5V/1.25V
Output Reference Levels	1.5V/1.25V
Output Load	Figure 1

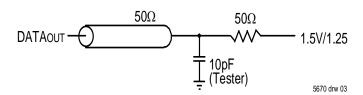


Figure 1. AC Output Test load.

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range⁽⁵⁾

			3/1S10 I Only	70T633/1S12 Com'l & Ind		70T633/1S15 Com'l Only		
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
READ CYCLE								
trc	Read Cycle Time	10		12		15		ns
t AA	Address Access Time		10		12		15	ns
tACE	Chip Enable Access Time ⁽³⁾		10		12		15	ns
t ABE	Byte Enable Access Time ⁽³⁾		5		6		7	ns
taoe	Output Enable Access Time		5		6		7	ns
tон	Output Hold from Address Change	3	_	3		3		ns
t_z	Output Low-Z Time ^(1,2)	0	_	0	_	0	_	ns
tHZ	Output High-Z Time ^(1,2)	0	4	0	6	0	8	ns
teu	Chip Enable to Power Up Time ⁽²⁾	0	_	0		0	_	ns
tpp	Chip Disable to Power Down Time ⁽²⁾		8		8	_	12	ns
tsop	Semaphore Flag Update Pulse (OE or SEM)		4		6		8	ns
tsaa	Semaphore Address Access Time	2	10	2	12	2	15	ns

5670 tb1 12

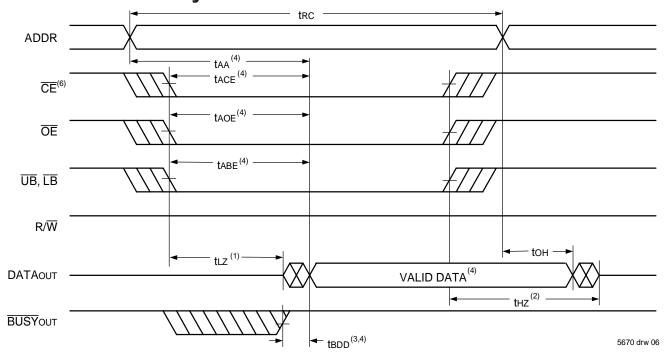
AC Electrical Characteristics Over the Operating Temperature and Supply Voltage⁽⁵⁾

		70T633/1S10 Com'l Only		70T633/1S12 Com'l & Ind		70T633/1S15 Com'l Only		
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
WRITE CYCLE								
twc	Write Cycle Time	10		12	_	15		ns
tew	Chip Enable to End-of-Write ⁽³⁾	8		10		12		ns
taw	Address Valid to End-of-Write	8		10		12		ns
tas	Address Set-up Time ⁽³⁾	0		0	_	0		ns
twp	Write Pulse Width	8		10		12		ns
twr	Write Recovery Time	0		0		0		ns
tow	Data Valid to End-of-Write	6		8		10		ns
tон	Data Hold Time ⁽⁴⁾	0		0		0		ns
twz	Write Enable to Output in High-Z ^(1,2)		4	_	6	_	8	ns
tow	Output Active from End-of-Write ^(1,2,4)	0		0		0		ns
tswrd	SEM Flag Write to Read Time	5		5		5		ns
tsps	SEM Flag Contention Window	5	_	5		5		ns

NOTES:

- 1. Transition is measured 0mV from Low or High-impedance voltage with Output Test Load (Figure 1).
- 2. This parameter is guaranted by device characterization, but is not production tested.
- 3. To access RAM, OE = VIL and SEM = VIH. To access semaphore, CE = VIH and SEM = VIL. Either condition must be valid for the entire tew time. CE = VIL when $\overline{\text{CE}}_0$ = V_IL and CE₁ = V_IH. $\overline{\text{CE}}$ = V_IH when $\overline{\text{CE}}_0$ = V_IH and/or CE₁ = V_IL.
- 4. The specification for toh must be met by the device supplying write data to the RAM under all operating conditions. Although toh and tow values will vary over voltage and temperature, the actual tDH will always be smaller than the actual tow.
- 5. These values are valid regardless of the power supply level selected for I/O and control signals (3.3V/2.5V). See page 6 for details.

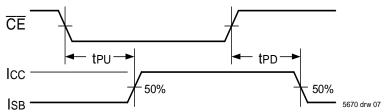
Waveform of Read Cycles⁽⁵⁾



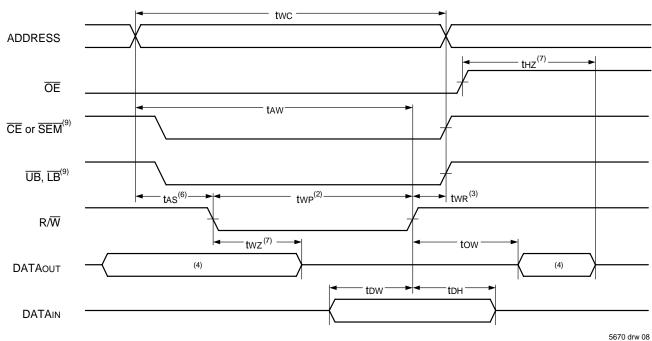
NOTES:

- 1. Timing depends on which signal is asserted last, \overline{OE} , \overline{CE} , \overline{LB} or \overline{UB} .
- 2. Timing depends on which signal is de-asserted first \overline{CE} , \overline{OE} , \overline{LB} or \overline{UB} .
- 3. tedd delay is required only in cases where the opposite port is completing a write operation to the same address location. For simultaneous read operations BUSY has no relation to valid output data.
- 4. Start of valid data depends on which timing becomes effective last: tAOE, tACE, tAA, tABE, or tBDD.
- 5. $\overline{SEM} = VIH.$
- 6. \overline{CE} = L occurs when \overline{CE} 0 = V_{IL} and CE₁ = V_{IH}. \overline{CE} = H when \overline{CE} 0 = V_{IH} and/or CE₁ = V_{IL}.

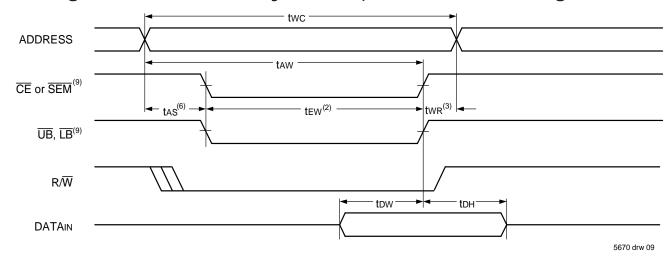
Timing of Power-Up Power-Down



Timing Waveform of Write Cycle No. 1, R/W Controlled Timing^(1,5,8)

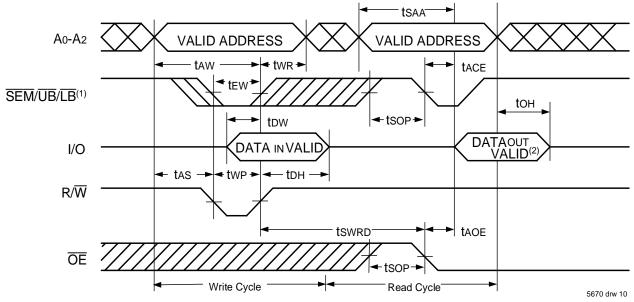


Timing Waveform of Write Cycle No. 2, $\overline{\text{CE}}$ Controlled Timing^(1,5,8)



- 1. $R\overline{W}$ or \overline{CE} or \overline{UB} or \overline{LB} = VIH during all address transitions.
- 3. twn is measured from the earlier of $\overline{\text{CE}}$ or $\overline{\text{R/W}}$ (or $\overline{\text{SEM}}$ or $\overline{\text{R/W}}$) going HIGH to the end of write cycle.
- 4. During this period, the I/O pins are in the output state and input signals must not be applied.
- 5. If the CE or SEM = ViL transition occurs simultaneously with or after the R/W = ViL transition, the outputs remain in the High-impedance state.
- 6. Timing depends on which enable signal is asserted last, $\overline{\text{CE}}$ or R/\overline{W} .
- 7. This parameter is guaranteed by device characterization, but is not production tested. Transition is measured 0mV from steady state with the Output Test Load (Figure 1).
- 8. If $\overline{OE} = \dot{V}_{IL}$ during R/W controlled write cycle, the write pulse width must be the larger of twp or (twz + tow) to allow the I/O drivers to turn off and data to be placed on the bus for the required tow. If $\overline{OE} = V_{IH}$ during an R/W controlled write cycle, this requirement does not apply and the write pulse can be as short as the specified twp.
- 9. To access RAM, $\overline{CE} = VIL$ and $\overline{SEM} = VIH$. To access semaphore, $\overline{CE} = VIH$ and $\overline{SEM} = VIL$. tew must be met for either condition. $\overline{CE} = VIL$ when $\overline{CE}_0 = VIL$ and $\overline{CE}_1 = VIL$.

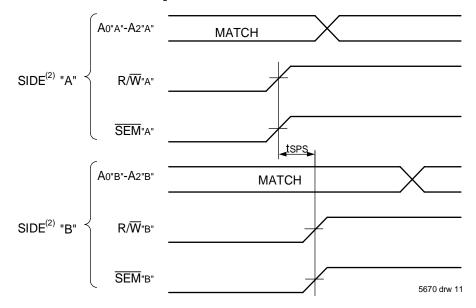
Timing Waveform of Semaphore Read after Write Timing, Either Side⁽¹⁾



NOTES:

- 1. $\overline{\text{CE}}_0 = \text{V}_{\text{IH}}$ and $\text{CE}_1 = \text{V}_{\text{IL}}$ are required for the duration of both the write cycle and the read cycle waveforms shown above. Refer to Truth Table II for details and for appropriate $\overline{\text{UB}}/\overline{\text{LB}}$ controls.
- 2. "DATAOUT VALID" represents all I/O's (I/Oo I/O17) equal to the semaphore value.

Timing Waveform of Semaphore Write Contention (1,3,4)



- 1. DOR = DOL = VIL, CEOL = CEOR = VIH; CE1L = CE1R = VIL. Refer also to Truth Table II for appropriate UB/LB controls.
- 2. All timing is the same for left and right ports. Port "A" may be either left or right port. "B" is the opposite from port "A".
- 3. This parameter is measured from R/W̄"a" or SEM̄"a" going HIGH to R/W̄"b" or SEM̄"b" going HIGH.
- 4. If tsps is not satisfied, the semaphore will fall positively to one side or the other, but there is no guarantee which side will be granted the semaphore flag.

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range

Ob-sl	Danmata		70T633/1S10 Com'l Only		70T633/1S12 Com'l & Ind		70T633/1S15 Com'l	
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
BUSY TIMIN	G (M/S=VIH)	•		•		•	•	
tbaa	BUSY Access Time from Address Match		10	_	12	_	15	ns
tbda	BUSY Disable Time from Address Not Matched	_	10		12	_	15	ns
tBAC	BUSY Access Time from Chip Enable Low	_	10		12		15	ns
tBDC	BUSY Disable Time from Chip Enable High	_	10		12		15	ns
taps	Arbitration Priority Set-up Time ⁽²⁾	2.5		2.5		2.5		ns
tBDD	BUSY Disable to Valid Data ⁽³⁾	_	10		12	_	15	ns
twн	Write Hold After BUSY ⁽⁵⁾	8		10		12		ns
BUSY TIMIN	G (M/S=VIL)							
twB	BUSY Input to Write ⁽⁴⁾	0	_	0		0		ns
twн	Write Hold After BUSY ⁽⁵⁾	8	_	10		12	_	ns
PORT-TO-PO	RT DELAY TIMING							
twdd	Write Pulse to Data Delay ⁽¹⁾	_	22		25		30	ns
todo	Write Data Valid to Read Data Delay(1)		20		22	_	25	ns

5670 tbl 14

NOTES:

- 1. Port-to-port delay through RAM cells from writing port to reading port, refer to "Timing Waveform of Write with Port-to-Port Read and BUSY (M/S = VIH)".
- 2. To ensure that the earlier of the two ports wins.
- 3. tBDD is a calculated parameter and is the greater of the Max. spec, twDD twp (actual), or tDDD tDw (actual).
- 4. To ensure that the write cycle is inhibited on port "B" during contention on port "A".
- 5. To ensure that a write cycle is completed on port "B" after contention on port "A".

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range^(1,2,3)

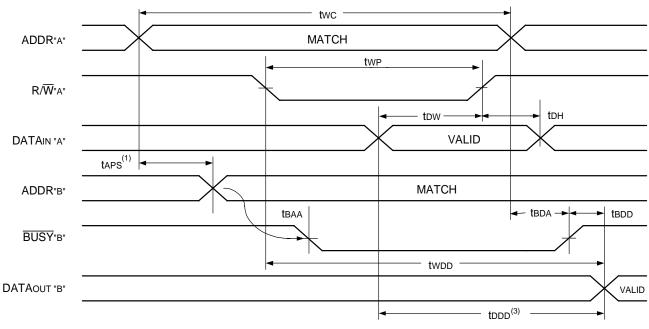
Cumhal	Dovementer	70T633/1S10 Com'l Only		70T633/1S12 Com'l & Ind		70T633/1S15 Com'l Only		
Symbol	Parameter		Max.	Min.	Max.	Min.	Max.	Unit
SLEEP MODE TIMING (ZZx=Vih)								
tzzs	Sleep Mode Set Time	10	_	12	_	15	_	ns
tzzr	Sleep Mode Reset Time	10	_	12	_	15	_	ns
tzzpd	Sleep Mode Power Down Time	10	_	12	_	15	_	ns
tzzpu	Sleep Mode Power Up Time	_	0		0	_	0	ns

NOTES:1. Timing is the same for both ports.

5670 tbl 14a

- 2. The sleep mode pin shuts off all dynamic inputs, except JTAG inputs, when asserted. OPTx, BUSYx, INTx, M/S and the sleep mode pins themselves (ZZx) are not affected during sleep mode. JTAG cannot be tested in sleep mode.
- 3. These values are valid regardless of the power supply level selected for I/O and control signals (3.3V/2.5V). See page 6 for details.

Timing Waveform of Write with Port-to-Port Read and \overline{BUSY} (M/ \overline{S} = VIH) $^{(2,4,5)}$

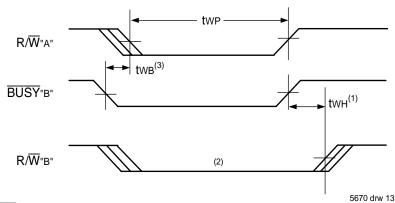


NOTES:

5670 drw 12

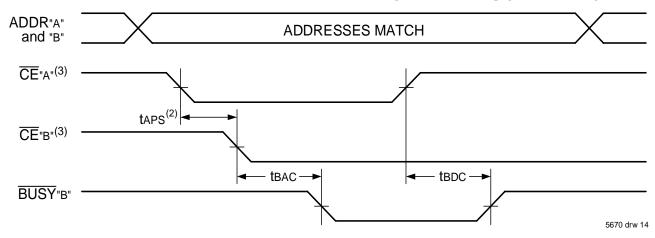
- 1. To ensure that the earlier of the two ports wins. taps is ignored for $M/\overline{S} = V_{IL}$ (SLAVE).
- 2. $\overline{\text{CE}}_{0L} = \overline{\text{CE}}_{0R} = \text{Vil}; \text{CE}_{1L} = \text{CE}_{1R} = \text{Vih}.$
- 3. \overline{OE} = VIL for the reading port.
- 4. If $M/\overline{S} = VIL$ (slave), \overline{BUSY} is an input. Then for this example $\overline{BUSY}^*A^* = VIH$ and \overline{BUSY}^*B^* input is shown above.
- 5. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".

Timing Waveform of Write with \overline{BUSY} (M/ \overline{S} = VIL)

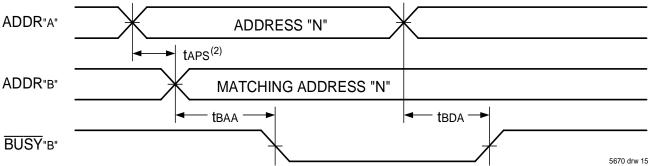


- NOTES
- 1. twh must be met for both $\overline{\text{BUSY}}$ input (SLAVE) and output (MASTER).
- 2. BUSY is asserted on port "B" blocking R/W"B", until BUSY"B" goes HIGH.
- 3. two only applies to the slave mode.

Waveform of BUSY Arbitration Controlled by CE Timing (M/S = Vih)(¹)



Waveform of \overline{BUSY} Arbitration Cycle Controlled by Address Match Timing (M/ \overline{S} = VIH)^(1,3,4)



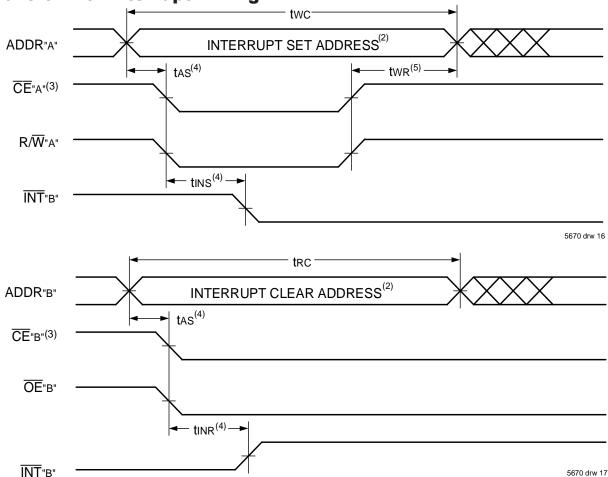
NOTES

- 1. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".
- 2. If taps is not satisfied, the BUSY signal will be asserted on one side or another but there is no guarantee on which side BUSY will be asserted.
- 3. $\overline{CE}x = V_{IL}$ when $\overline{CE}_{0x} = V_{IL}$ and $\overline{CE}_{1x} = V_{IH}$. $\overline{CE}x = V_{IH}$ when $\overline{CE}_{0x} = V_{IH}$ and/or $\overline{CE}_{1x} = V_{IL}$.
- 4. $\overline{CE}_{0x} = \overline{OE}_{x} = \overline{LB}_{x} = \overline{UB}_{x} = V_{IL}$. $CE_{1x} = V_{IH}$.

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range

			3/1S10 'I Only	Co	3/1S12 70T633/1S15 om'l Com'l Only Ind			
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
INTERRUPT :	TIMING							
tas	Address Set-up Time	0	_	0	—	0	_	ns
twr	Write Recovery Time	0	_	0	_	0	_	ns
tins	Interrupt Set Time	_	10	_	12	_	15	ns
tinr	Interrupt Reset Time		10		12		15	ns

Waveform of Interrupt Timing⁽¹⁾



NOTES:

- 1. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".
- 2. Refer to Interrupt Truth Table.
- 3. $\overline{CE}x = V_{IL}$ means $\overline{CE}_{0x} = V_{IL}$ and $CE_{1x} = V_{IH}$. $\overline{CE}x = V_{IH}$ means $\overline{CE}_{0x} = V_{IH}$ and/or $CE_{1x} = V_{IL}$.
- 4. Timing depends on which enable signal $(\overline{CE} \text{ or } R/\overline{W})$ is asserted last.
- 5. Timing depends on which enable signal (CE or R/W) is de-asserted first.

Truth Table III — Interrupt Flag^(1,4)

Left Port					Right Port					
R/W̄∟	CEL	ŌĒL	A18L-A0L ⁽⁵⁾	ĪNT∟	R/W̄R	CER	Ō E R	A18R-A0R ⁽⁵⁾	ĪNT⊓	Function
L	L	Х	7FFFF	Х	Х	Х	Х	Х	L ⁽²⁾	Set Right INTR Flag
Х	Х	Х	Х	Х	Х	L	L	7FFFF	H ⁽³⁾	Reset Right INTR Flag
Х	Х	Х	Х	L ⁽³⁾	L	L	Х	7FFFE	Х	Set Left INTL Flag
Х	L	L	7FFFE	H ⁽²⁾	Х	Х	Х	Х	Х	Reset Left INTL Flag

NOTES:

- 1. Assumes $\overline{BUSY}L = \overline{BUSY}R = VIH$. $\overline{CE}X = L$ means $\overline{CE}0X = VIL$ and CE1X = VIH.
- 2. If $\overline{BUSY}L = VIL$, then no change.
- 3. If $\overline{BUSYR} = VIL$, then no change.
- 4. $\overline{\text{INT}}_{\text{L}}$ and $\overline{\text{INT}}_{\text{R}}$ must be initialized at power-up.
- 5. A18x is a NC for IDT70T631. Therefore, Interrupt Addresses are 3FFFF and 3FFFE.

Truth Table IV — Address BUSY Arbitration

	In	puts	Out	puts	
CEL ⁽⁵⁾	CER ⁽⁵⁾	A0L-A18L ⁽⁴⁾ A0R-A18R	BUSY _L (1)	BUSY _R (1)	Function
Х	Χ	NO MATCH	Н	Н	Normal
Н	Χ	MATCH	Н	Н	Normal
Х	Н	MATCH	Н	Н	Normal
L	L	MATCH	(2)	(2)	Write Inhibit ⁽³⁾

5670 tbl 17

NOTES:

- 1. Pins BUSYL and BUSYR are both outputs when the part is configured as a master. Both are inputs when configured as a slave. BUSY outputs on the IDT70T633/1 are push-pull, not open drain outputs. On slaves the BUSY input internally inhibits writes.
- 2. "L" if the inputs to the opposite port were stable prior to the address and enable inputs of this port. "H" if the inputs to the opposite port became stable after the address and enable inputs of this port. If taps is not met, either BUSYL or BUSYR = LOW will result. BUSYL and BUSYR outputs can not be LOW simultaneously.
- 3. Writes to the left port are internally ignored when BUSYL outputs are driving LOW regardless of actual logic level on the pin. Writes to the right port are internally ignored when BUSYR outputs are driving LOW regardless of actual logic level on the pin.
- 4. A18 is a NC for IDT70T631. Address comparison will be for A0 A17.
- 5. $\overline{CE}x = L$ means $\overline{CE}0x = VIL$ and CE1x = VIH. $\overline{CE}x = H$ means $\overline{CE}0x = VIH$ and/or CE1x = VIL.

Truth Table V — Example of Semaphore Procurement Sequence (1,2,3)

HACH LABIO L ZXA	ilipio di de	THE PHOTO	i roodronioni ooquonoo
Functions	Do - D17 Left	Do - D17 Right	Status
No Action	1	1	Semaphore free
Left Port Writes "0" to Semaphore	0	1	Left port has semaphore token
Right Port Writes "0" to Semaphore	0	1	No change. Right side has no write access to semaphore
Left Port Writes "1" to Semaphore	1	0	Right port obtains semaphore token
Left Port Writes "0" to Semaphore	1	0	No change. Left port has no write access to semaphore
Right Port Writes "1" to Semaphore	0	1	Left port obtains semaphore token
Left Port Writes "1" to Semaphore	1	1	Semaphore free
Right Port Writes "0" to Semaphore	1	0	Right port has semaphore token
Right Port Writes "1" to Semaphore	1	1	Semaphore free
Left Port Writes "0" to Semaphore	0	1	Left port has semaphore token
Left Port Writes "1" to Semaphore	1	1	Semaphore free

NOTES: 5670 tbl 18

- 1. This table denotes a sequence of events for only one of the eight semaphores on the IDT70T633/1.
- 2. There are eight semaphore flags written to via I/Oo and read from all I/O's (I/Oo-I/O17). These eight semaphores are addressed by Ao A2.
- 3. CEo = VIH, CE1 = SEM = VIL to access the semaphores. Refer to the Semaphore Read/Write Control Truth Table.

Functional Description

The IDT70T633/1 provides two ports with separate control, address and I/O pins that permit independent access for reads or writes to any location in memory. The IDT70T633/1 has an automatic power down feature controlled by $\overline{\text{CE}}$. The $\overline{\text{CE}}$ 0 and CE1 control the on-chip power down circuitry that permits the respective port to go into a standby mode when not selected ($\overline{\text{CE}}$ = HIGH). When a port is enabled, access to the entire memory array is permitted.

Interrupts

If the user chooses the interrupt function, a memory location (mail box or message center) is assigned to each port. The left port interrupt

flag ($\overline{\text{INTL}}$) is asserted when the right port writes to memory location 7FFFE (HEX), where a write is defined as $\overline{\text{CE}}_R = R/\overline{W}_R = V_{IL}$ per the Truth Table. The left port clears the interrupt through access of address location 7FFFE when $\overline{\text{CE}}_L = \overline{\text{OE}}_L = V_{IL}$, R/\overline{W} is a "don't care". Likewise, the right port interrupt flag ($\overline{\text{INTR}}$) is asserted when the left port writes to memory location 7FFFF (HEX) and to clear the interrupt flag ($\overline{\text{INTR}}$), the right port must read the memory location 7FFFF. The message (18 bits) at 7FFFE or 7FFFF (3FFFF or 3FFFE for IDT70T631) is user-defined since it is an addressable SRAM location. If the interrupt function is not used, address locations 7FFFE and 7FFFF are not used as mail boxes, but as part of the random access memory. Refer to Truth Table III for the interrupt operation.

Preliminary

High-Speed 2.5V 512/256K x 18 Asynchronous Dual-Port Static RAM

Busy Logic

Busy Logic provides a hardware indication that both ports of the RAM have accessed the same location at the same time. It also allows one of the two accesses to proceed and signals the other side that the RAM is "Busy". The $\overline{\text{BUSY}}$ pin can then be used to stall the access until the operation on the other side is completed. If a write operation has been attempted from the side that receives a $\overline{\text{BUSY}}$ indication, the write signal is gated internally to prevent the write from proceeding.

The use of \overline{BUSY} logic is not required or desirable for all applications. In some cases it may be useful to logically OR the \overline{BUSY} outputs together and use any \overline{BUSY} indication as an interrupt source to flag the event of an illegal or illogical operation. If the write inhibit function of \overline{BUSY} logic is not desirable, the \overline{BUSY} logic can be disabled by placing the part in slave mode with the M/\overline{S} pin. Once in slave mode the \overline{BUSY} pin operates solely as a write inhibit input pin. Normal operation can be programmed by tying the \overline{BUSY} pins HIGH. If desired, unintended write operations can be prevented to a port by tying the \overline{BUSY} pin for that port LOW.

The BUSY outputs on the IDT70T633/1 RAM in master mode, are push-pull type outputs and do not require pull up resistors to operate. If these RAMs are being expanded in depth, then the BUSY indication for the resulting array requires the use of an external AND gate.

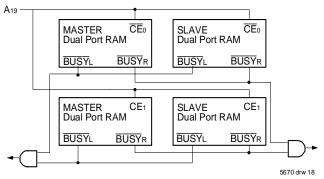


Figure 3. Busy and chip enable routing for both width and depth expansion with IDT70T633/1 Dual-Port RAMs.

Width Expansion with Busy Logic Master/Slave Arrays

When expanding an IDT70T633/1 RAM array in width while using $\overline{\text{BUSY}}$ logic, one master part is used to decide which side of the RAMs array will receive a $\overline{\text{BUSY}}$ indication, and to output that indication. Any number of slaves to be addressed in the same address range as the master use the $\overline{\text{BUSY}}$ signal as a write inhibit signal. Thus on the IDT70T633/1 RAM the $\overline{\text{BUSY}}$ pin is an output if the part is used as a master (M/ $\overline{\text{S}}$ pin = VIH), and the $\overline{\text{BUSY}}$ pin is an input if the part used as a slave (M/ $\overline{\text{S}}$ pin = VIL) as shown in Figure 3.

If two or more master parts were used when expanding in width, a split decision could result with one master indicating BUSY on one side of the array and another master indicating BUSY on one other side of the array. This would inhibit the write operations from one port for part of a word and inhibit the write operations from the other port for the other part of the word.

The BUSY arbitration on a master is based on the chip enable and

address signals only. It ignores whether an access is a read or write. In a master/slave array, both address and chip enable must be valid long enough for a $\overline{\text{BUSY}}$ flag to be output from the master before the actual write pulse can be initiated with the R/\overline{W} signal. Failure to observe this timing can result in a glitched internal write inhibit signal and corrupted data in the slave.

Semaphores

The IDT70T633/1 is an extremely fast Dual-Port $512/256K \times 18$ CMOS Static RAM with an additional 8 address locations dedicated to binary semaphore flags. These flags allow either processor on the left or right side of the Dual-Port RAM to claim a privilege over the other processor for functions defined by the system designer's software. As an example, the semaphore can be used by one processor to inhibit the other from accessing a portion of the Dual-Port RAM or any other shared resource.

The Dual-Port RAM features a fast access time, with both ports being completely independent of each other. This means that the activity on the left port in no way slows the access time of the right port. Both ports are identical in function to standard CMOS Static RAM and can be read from or written to at the same time with the only possible conflict arising from the simultaneous writing of, or a simultaneous READ/WRITE of, a non-semaphore location. Semaphores are protected against such ambiguous situations and may be used by the system program to avoid any conflicts in the non-semaphore portion of the Dual-Port RAM. These devices have an automatic power-down feature controlled by $\overline{\text{CE}}$ and $\overline{\text{CE}}$ 0, CE1, and $\overline{\text{SEM}}$ pins control onchip power down circuitry that permits the respective port to go into standby mode when not selected.

Systems which can best use the IDT70T633/1 contain multiple processors or controllers and are typically very high-speed systems which are software controlled or software intensive. These systems can benefit from a performance increase offered by the hardware semaphores of the IDT70T633/1, which provide a lockout mechanism without requiring complex programming.

Software handshaking between processors offers the maximum in system flexibility by permitting shared resources to be allocated in varying configurations. The IDT70T633/1 does not use its semaphore flags to control any resources through hardware, thus allowing the system designer total flexibility in system architecture.

An advantage of using semaphores rather than the more common methods of hardware arbitration is that wait states are never incurred in either processor. This can prove to be a major advantage in very high-speed systems.

How the Semaphore Flags Work

The semaphore logic is a set of eight latches which are independent of the Dual-Port RAM. These latches can be used to pass a flag, or token, from one port to the other to indicate that a shared resource is in use. The semaphores provide a hardware assist for a use assignment method called "Token Passing Allocation." In this method, the state of a semaphore latch is used as a token indicating that a shared resource is in use. If the left processor wants to use this resource, it requests the token by setting the latch. This processor then

verifies its success in setting the latch by reading it. If it was successful, it proceeds to assume control over the shared resource. If it was not successful in setting the latch, it determines that the right side processor has set the latch first, has the token and is using the shared resource. The left processor can then either repeatedly request that semaphore's status or remove its request for that semaphore to perform another task and occasionally attempt again to gain control of the token via the set and test sequence. Once the right side has relinguished the token, the left side should succeed in gaining control.

The semaphore flags are active LOW. A token is requested by writing a zero into a semaphore latch and is released when the same side writes a one to that latch.

The eight semaphore flags reside within the IDT70T633/1 in a separate memory space from the Dual-Port RAM array. This address space is accessed by placing a low input on the SEM pin (which acts as a chip select for the semaphore flags) and using the other control pins (Address, \overline{CE}_0 , \overline{CE}_1 , $\overline{R/W}$ and $\overline{LB}/\overline{UB}$) as they would be used in accessing a standard Static RAM. Each of the flags has a unique address which can be accessed by either side through address pins A0-A2. When accessing the semaphores, none of the other address pins has any effect.

When writing to a semaphore, only data pin Do is used. If a low level is written into an unused semaphore location, that flag will be set to a zero on that side and a one on the other side (see Truth Table V). That semaphore can now only be modified by the side showing the zero. When a one is written into the same location from the same side, the flag will be set to a one for both sides (unless a semaphore request from the other side is pending) and then can be written to by both sides. The fact that the side which is able to write a zero into a semaphore subsequently locks out writes from the other side is what makes semaphore flags useful in interprocessor communications. (A thorough discussion on the use of this feature follows shortly.) A zero written into the same location from the other side will be stored in the semaphore request latch for that side until the semaphore is freed by the first side.

When a semaphore flag is read, its value is spread into all data bits so that a flag that is a one reads as a one in all data bits and a flag containing a zero reads as all zeros. The read value is latched into one side's output register when that side's semaphore, byte select (SEM, $\overline{LB}/\overline{UB}$) and output enable (\overline{OE}) signals go active. This serves to disallow the semaphore from changing state in the middle of a read cycle due to a write cycle from the other side.

A sequence WRITE/READ must be used by the semaphore in order to guarantee that no system level contention will occur. A processor requests access to shared resources by attempting to write a zero into a semaphore location. If the semaphore is already in use, the semaphore request latch will contain a zero, yet the semaphore flag will appear as one, a fact which the processor will verify by the subsequent read (see Table V). As an example, assume a processor writes a zero to the left port at a free semaphore location. On a subsequent read, the processor will verify that it has written successfully to that location and will assume control over the resource in question. Meanwhile, if a processor on the right side attempts to write a zero to the same semaphore flag it will fail, as will be verified by the fact that a one will be read from that semaphore on the right side during subsequent read. Had a sequence of READ/WRITE been used instead, system contention problems could have occurred during

the gap between the read and write cycles.

It is important to note that a failed semaphore request must be followed by either repeated reads or by writing a one into the same location. The reason for this is easily understood by looking at the simple logic diagram of the semaphore flag in Figure 4. Two semaphore request latches feed into a semaphore flag. Whichever latch is first to present a zero to the semaphore flag will force its side of the semaphore flag LOW and the opposite side HIGH. This condition will continue until a one is written to the same semaphore request latch. If the opposite side semaphore request latch has been written to zero in the meantime, the semaphore flag will flip over to the other side as soon as a one is written into the first request latch. The opposite side flag will now stay LOW until its semaphore request latch

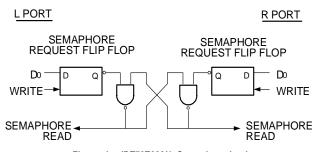


Figure 4. IDT70T633/1 Semaphore Logic 5670 drw 19

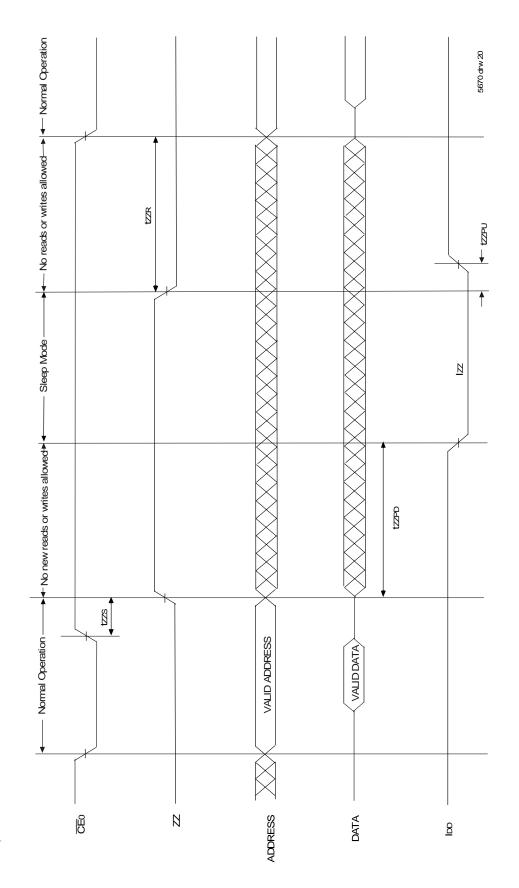
is written to a one. From this it is easy to understand that, if a semaphore is requested and the processor which requested it no longer needs the resource, the entire system can hang up until a one is written into that semaphore request latch.

The critical case of semaphore timing is when both sides request a single token by attempting to write a zero into it at the same time. The semaphore logic is specially designed to resolve this problem. If simultaneous requests are made, the logic guarantees that only one side receives the token. If one side is earlier than the other in making the request, the first side to make the request will receive the token. If both requests arrive at the same time, the assignment will be arbitrarily made to one port or the other.

One caution that should be noted when using semaphores is that semaphores alone do not guarantee that access to a resource is secure. As with any powerful programming technique, if semaphores are misused or misinterpreted, a software error can easily happen.

Initialization of the semaphores is not automatic and must be handled via the initialization program at power-up. Since any semaphore request flag which contains a zero must be reset to a one, all semaphores on both sides should have a one written into them at initialization from both sides to assure that they will be free when needed.

Timing Waveform of Sleep Mode^(1,2)



Sleep Mode

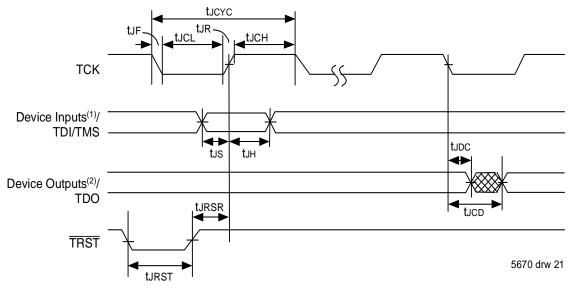
The IDT70T633/1 is equipped with an optional sleep or low power mode on both ports. The sleep mode pin on both ports is active high. During normal operation, the ZZ pin is pulled low. When ZZ is pulled high, the port will enter sleep mode where it will have the lowest possible power consumption. The sleep mode timing diagram demonstrates the modes of operation: Normal Operation, No Read/Write Allowed and Sleep Mode.

For a period of time prior to sleep mode and after recovering from sleep

mode (tzzs and tzzr), new reads or writes are not allowed. If a write or read operation occurs during these periods, the memory array may be corrupted. Validity of data out from the RAM cannot be guaranteed immediately after ZZ is asserted (prior to being in sleep).

During sleep mode the RAM automatically deselects itself and disconnects its internal buffer. All outputs will remain in high-Z state while in sleep mode. All inputs are allowed to toggle, but the RAM will not be selected and will not perform any reads or writes.

JTAG Timing Specifications



NOTES:

- 1. Device inputs = All device inputs except TDI, TMS, and TRST.
- 2. Device outputs = All device outputs except TDO.

JTAG AC Electrical Characteristics^(1,2,3,4,5)

			70T633/1	
Symbol	Parameter	Min.	Max.	Units
tucyc	JTAG Clock Input Period	100		ns
tлсн	JTAG Clock HIGH	40	_	ns
tucL	JTAG Clock Low	40	_	ns
tur	JTAG Clock Rise Time	_	3 ⁽¹⁾	ns
₩F	JTAG Clock Fall Time	_	3 ⁽¹⁾	ns
tJRST	JTAG Reset	50	_	ns
tursr	JTAG Reset Recovery	50		ns
tico	JTAG Data Output	_	25	ns
tudo	JTAG Data Output Hold	0		ns
tus	JTAG Setup	15		ns
tлн	JTAG Hold	15		ns

5670 tbl 19

- 1. Guaranteed by design.
- 2. 30pF loading on external output signals.
- 3. Refer to AC Electrical Test Conditions stated earlier in this document.
- JTAG operations occur at one speed (10MHz). The base device may run at any speed specified in this datasheet.
- 5. JTAG cannot be tested in sleep mode.

Identification Register Definitions

Instruction Field	Value	Description
Revision Number (31:28)	0x0	Reserved for version number
IDT Device ID (27:12)	0x33B ⁽¹⁾	Defines IDT part number 70T633
IDT JEDEC ID (11:1)	0x33	Allows unique identification of device vendor as IDT
ID Register Indicator Bit (Bit 0)	1	Indicates the presence of an ID register

NOTE:

5670 tbl 20

Scan Register Sizes

Register Name	Bit Size
Instruction (IR)	4
Bypass (BYR)	1
Identification (IDR)	32
Boundary Scan (BSR)	Note (3)

5670 tbl 21

System Interface Parameters

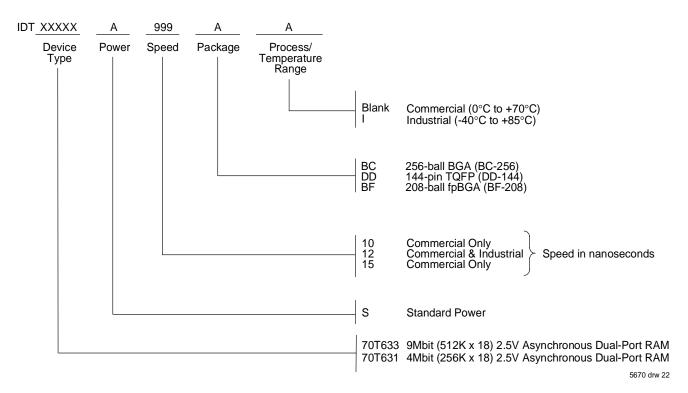
Instruction	Code	Description
EXTEST	0000	Forces contents of the boundary scan cells onto the device outputs ⁽¹⁾ . Places the boundary scan register (BSR) between TDI and TDO.
BYPASS	1111	Places the bypass register (BYR) between TDI and TDO.
IDCODE	0010	Loads the ID register (IDR) with the vendor ID code and places the register between TDI and TDO.
HIGHZ	0100	Places the bypass register (BYR) between TDI and TDO. Forces all device output drivers to a High-Z state.
CLAMP	0011	Uses BYR. Forces contents of the boundary scan cells onto the device outputs. Places the bypass register (BYR) between TDI and TDO.
SAMPLE/PRELOAD	0001	Places the boundary scan register (BSR) between TDI and TDO. SAMPLE allows data from device inputs ^[2] and outputs ⁽¹⁾ to be captured in the boundary scan cells and shifted serially through TDO. PRELOAD allows data to be input serially into the boundary scan cells via the TDI.
RESERVED	All other codes	Several combinations are reserved. Do not use codes other than those identified above.

NOTES

- 1. Device outputs = All device outputs except TDO.
- 2. Device inputs = All device inputs except TDI, TMS, and $\overline{\text{TRST}}$.
- 3. The Boundary Scan Descriptive Language (BSDL) file for this device is available on the IDT website (www.idt.com), or by contacting your local IDT sales representative.

^{1.} Device ID for IDT70T631 is 0x33C.

Ordering Information



Preliminary Datasheet: Definition

"PRELIMINARY' datasheets contain descriptions for products that are in early release.

Datasheet Document History:

04/25/03: Initial Datasheet



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